

### Theorem 2.1.1 Logical Equivalences

Given any statement variables  $p$ ,  $q$ , and  $r$ , a tautology  $\mathbf{t}$  and a contradiction  $\mathbf{c}$ , the following logical equivalences hold.

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| 1. Commutative laws: $p \wedge q \equiv q \wedge p$                                  | $p \vee q \equiv q \vee p$                                |
| 2. Associative laws: $(p \wedge q) \wedge r \equiv p \wedge (q \wedge r)$            | $(p \vee q) \vee r \equiv p \vee (q \vee r)$              |
| 3. Distributive laws: $p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$    | $p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$ |
| 4. Identity laws: $p \wedge \mathbf{t} \equiv p$                                     | $p \vee \mathbf{c} \equiv p$                              |
| 5. Negation laws: $p \vee \sim p \equiv \mathbf{t}$                                  | $p \wedge \sim p \equiv \mathbf{c}$                       |
| 6. Double negative law: $\sim(\sim p) \equiv p$                                      |   |
| 7. Idempotent laws: $p \wedge p \equiv p$  | $p \vee p \equiv p$                                       |
| 8. Universal bound laws: $p \vee \mathbf{t} \equiv \mathbf{t}$                       | $p \wedge \mathbf{c} \equiv \mathbf{c}$                   |
| 9. De Morgan's laws: $\sim(p \wedge q) \equiv \sim p \vee \sim q$                    | $\sim(p \vee q) \equiv \sim p \wedge \sim q$              |
| 10. Absorption laws: $p \vee (p \wedge q) \equiv p$                                  | $p \wedge (p \vee q) \equiv p$                            |
| 11. Negations of $\mathbf{t}$ and $\mathbf{c}$ : $\sim \mathbf{t} \equiv \mathbf{c}$ | $\sim \mathbf{c} \equiv \mathbf{t}$                       |

**Table 2.3.1 Valid Argument Forms (Rules of Inference)**

<b>Modus Ponens</b>	$p \rightarrow q$ $p$ $\therefore q$	<b>Elimination</b>	a. $p \vee q$ $\sim q$ $\therefore p$	b. $p \vee q$ $\sim p$ $\therefore q$
<b>Modus Tollens</b>	$p \rightarrow q$ $\sim q$ $\therefore \sim p$	<b>Transitivity</b>	$p \rightarrow q$ $q \rightarrow r$ $\therefore p \rightarrow r$	
<b>Generalization</b>	a. $p$ $\therefore p \vee q$	b. $q$ $\therefore p \vee q$	<b>Proof by</b>	$p \vee q$
			<b>Division into Cases</b>	$p \rightarrow r$ $q \rightarrow r$ $\therefore r$
<b>Specialization</b>	a. $p \wedge q$ $\therefore p$	b. $p \wedge q$ $\therefore q$		
<b>Conjunction</b>	$p$ $q$ $\therefore p \wedge q$	<b>Contradiction Rule</b>	$\sim p \rightarrow \mathbf{c}$ $\therefore p$	